Algorithm 790: CSHEP2D: Cubic Shepard Method for Bivariate Interpolation of Scattered Data

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We describe a new algorithm for scattered data interpolation. The method is similar to that of Algorithm 660 but achieves cubic precision and C^2 continuity at very little additional cost. An accompanying article presents test results that show the method to be among the most accurate available.

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1. METHOD

We treat the problem of constructing a smooth bivariate function C that interpolates data values f_k at scattered nodes (x_k, y_k) in the plane for $k = 1, \ldots N$. We employ a modified Shepard method with a cell-based search algorithm as described in Renka [1988a]. The interpolant is defined by

$$C(x, y) = \sum_{k=1}^{N} W_k(x, y) C_k(x, y) \Big/ \sum_{i=1}^{N} W_i(x, y),$$

where the nodal function C_k is a bivariate cubic polynomial that interpolates the data value f_k at node k and fits the data values on a set of nearby nodes in a weighted least-squares sense.

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The unnormalized weights are inverse distance functions:

$$W_k(x, y) = \left[\frac{(R_w - d_k)_+}{R_w d_k} \right]^3$$

for

$$(R_w - d_k)_+ = \begin{cases} R_w - d_k & \text{if } d_k < R_w \\ 0 & \text{if } d_k \geq R_w \end{cases},$$

where $d_k(x, y)$ is the Euclidean distance between (x, y) and (x_k, y_k), and R_w is a radius of influence about (x_k, y_k). An interpolated value at a point (x, y) depends only on the data at nodes whose radii include (x, y).

It follows from the above definition that C interpolates the data, maintains the local shape properties of the nodal functions (has first and second partial derivatives at (x_k, y_k) that agree with those of C_k), has cubic precision, and lies in the space $C^2(\mathbf{R}^2)$.

Nodal function C_k is defined by

$$\begin{aligned} C_k(x, y) = {} & a_{1k}(x - x_k)^3 + a_{2k}(x - x_k)^2(y - y_k) + a_{3k}(x - x_k)(y - y_k)^2 \\ & + a_{4k}(y - y_k)^3 + a_{5k}(x - x_k)^2 + a_{6k}(x - x_k)(y - y_k) \\ & + a_{7k}(y - y_k)^2 + a_{8k}(x - x_k) + a_{9k}(y - y_k) + f_k, \end{aligned}$$

where the coefficients minimize

$$\sum_{\substack{i=1 \\ i \neq k}}^{N} \omega_{ik}[a_{1k}(x_i - x_k)^3 + \ldots + a_{9k}(y_i - y_k) + f_k - f_i]^2$$

for

$$\omega_{ik} = \left[\frac{(R_c - d_{ik})_+}{R_c d_{ik}} \right]^2,$$

where d_{ik} is the distance between nodes i and k, and R_c is a radius of influence about node k—C_k depends only on the data values at nodes within distance R_c of (x_k, y_k).

The radii R_c and R_w vary with k and are taken to be just large enough to include N_c and N_w nodes, respectively, for fixed values of N_c and N_w. The optimal values of these parameters depend on the data set, but accuracy varies smoothly and gradually with variations in the values. The default recommendations, found to be optimal for a set of test cases, are $N_c = 17$ and $N_w = 30$. However, for a nearly uniform rectangular grid of nodes, $N_c = 9$ was found to be optimal.

Note that, in general, the support of C is the union of a set of node-centered disks with radii that depend on N_w. If the nodal density varies widely, this union of disks may not cover the convex hull of the nodes, i.e., the convex hull could include points (x, y) for which $C(x, y) = 0$ because (x, y) is not within the radius of influence of any node. Thus, it may be necessary to use a larger value of N_w in order to avoid this situation.

The cell-based search method is used in the preprocessing phase to determine an ordered sequence of nearest neighbors to each node, and in the evaluation phase to find the set of all nodes whose radii R_w include the evaluation point. The smallest rectangle containing the nodes is partitioned into an $N_r \times N_r$ uniform grid of cells, and the indexes of the nodes contained in each cell are stored as a linked list in two integer arrays. Assuming a uniform distribution of nodes, the expected time complexity is $O(N)$ for the preprocessing phase and constant for each evaluation. Worst-case operation counts are $O(N^2)$ for preprocessing and $O(N)$ for evaluation.

The accompanying survey article [Renka and Brown 1999] presents test results showing, that for reasonably dense data sets, CSHEP2D is among the most accurate scattered data algorithms available.

2. CODE

The software is written in 1977 ANSI Standard Fortran and uses double precision. It can be converted to single precision (to save storage or to run on a Cray) by simply replacing all occurrences of 'DOUBLE PRECISION' or 'DBLE' by 'REAL'. There are no system dependencies. The array storage requirements consist of three order-N arrays, X, Y, F, containing the data points, a $9 \times N$ array A for the coefficients, an array RW of length N for the weights W_k, an $N_r \times N_r$ integer array LCELL for the index of the first node in each cell, and an integer array LNEXT of length N for next-node indexes.

The code is modularized in a fashion similar to that of Algorithm 660 [Renka 1988b]. The user-callable subprograms are as follows:

CSHEP2 Subroutine which computes the parameters defining the interpolant C.

CS2VAL Function which returns the value of C at an arbitrary point.

CS2GRD Subroutine which returns the value and gradient of C at an arbitrary point.

CS2HES Subroutine which returns the value, gradient, and Hessian of C at an arbitrary point.

STORE2 Subroutine which computes and stores the data structure for cell-based searches.

GETNP2 Subroutine which returns the nearest unmarked node, along with its Euclidean distance, to an arbitrary point, and marks the node (so that a subsequent call will return the next closest node).

CSHEP2 calls STORE2 and GETNP2 to find sequences of nearest neighbors to each node. It calls three additional subroutines to set up and solve the least-squares systems for the coefficients defining C. There are no other subprogram dependencies. Subroutines STORE2 and GETNP2 could be extracted from the source code and used to solve more general closest-point problems.

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